

软件分析与验证前沿

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Dynamic Symbolic Execution

Motivation

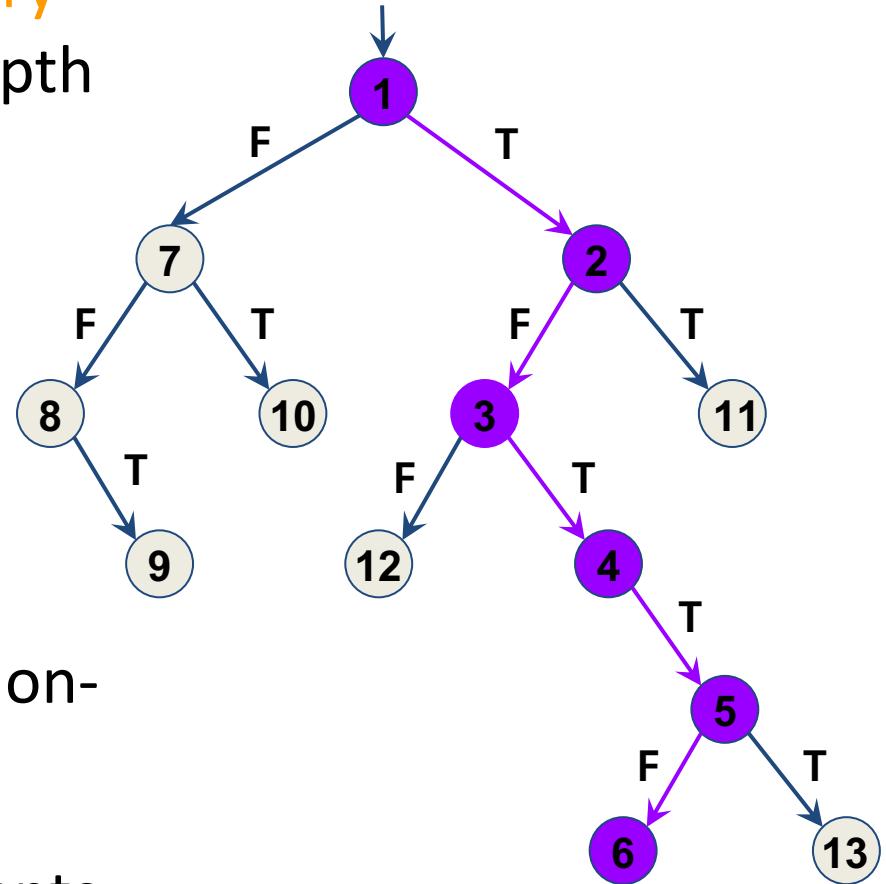
- Writing and maintaining tests is tedious and error-prone
- Idea: Automated Test Generation
 - Generate regression test suite
 - Execute all reachable statements
 - Catch any assertion violations

Approach

- Dynamic Symbolic Execution
 - Stores program state **concretely** and **symbolically**
 - Solves **constraints** to guide execution at branch points
 - Explores **all execution paths** of the unit tested
- Example of Hybrid Analysis
 - Collaboratively combines dynamic and static analysis

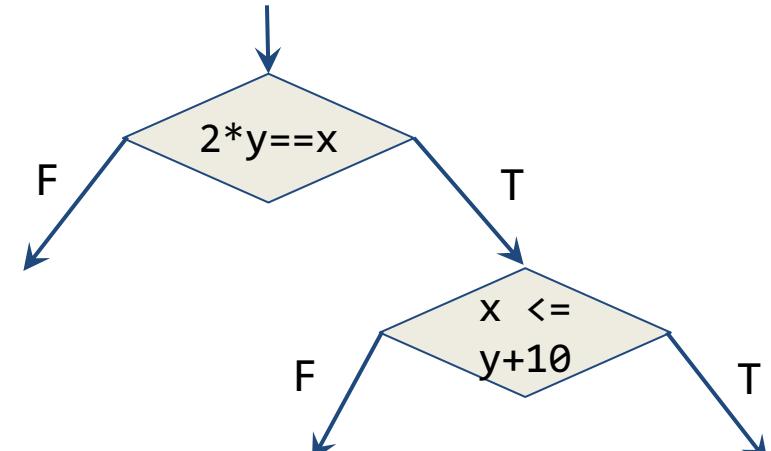
Execution Paths of a Program

- Program can be seen as **binary tree** with possibly infinite depth
 - Called **Computation Tree**
- Each **node** represents the execution of a conditional statement
- Each **edge** represents the execution of a sequence of non-conditional statements
- Each **path** in the tree represents an equivalence class of inputs



Example of Computation Tree

```
void test_me(int x, int y) {  
    if (2*y == x) {  
        if (x <= y+10)  
            print("OK");  
        else {  
            print("something bad");  
            ERROR;  
        }  
    } else  
        print("OK");  
}
```



```
assert(b) → if (!b) ERROR;
```

Existing Approach I

Random Testing

- Generate random inputs
- Execute the program on those (concrete) inputs

```
void test_me(int x) {  
    if (x == 94389) {  
        ERROR;  
    }  
}
```

Problem:

- Probability of reaching error could be astronomically small

Probability of **ERROR**:

$$1/2^{32} \approx 0.000000023\%$$

Existing Approach II

Symbolic Execution

- Use symbolic values for inputs
- Execute program symbolically on symbolic input values
- Collect symbolic path constraints
- Use **theorem prover** to check if a branch can be taken

```
void test_me(int x) {  
    if (x*3 == 15) {  
        if (x % 5 == 0)  
            print("OK");  
        else {  
            print("something bad");  
            ERROR;  
        }  
    } else  
        print("OK");  
}
```

Problem:

- Does not scale for large programs

Existing Approach II

Symbolic Execution

- Use symbolic values for inputs
- Execute program symbolically on symbolic input values
- Collect symbolic path constraints
- Use **theorem prover** to check if a branch can be taken

```
void test_me(int x) {  
    // c = product of two  
    // large primes  
    if (pow(2,x) % c == 17) {  
        print("something bad");  
        ERROR;  
    } else  
        print("OK");  
}
```

Symbolic execution will say both branches are reachable: **False Positive**

Problem:

- Does not scale for large programs

Combined Approach

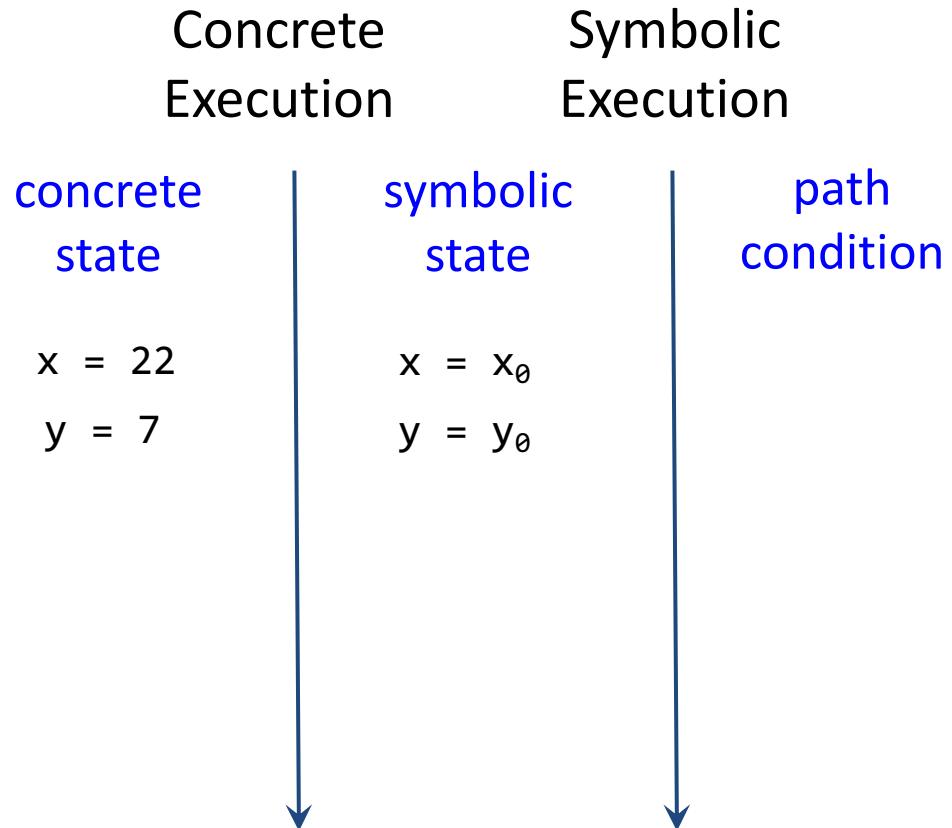
Dynamic Symbolic Execution (DSE)

- Start with random input values
- Keep track of **both** concrete values and symbolic constraints
- Use concrete values to **simplify** symbolic constraints
- **Incomplete** theorem-prover

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y); ←  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 22	x = x_0	
y = 7	y = y_0	
z = 14	z = $2*y_0$	

The diagram illustrates the flow of information from concrete execution to symbolic execution. It consists of three columns: 'Concrete Execution' (left), 'Symbolic Execution' (middle), and 'path condition' (right). Arrows point downwards from each column to the next. The 'Concrete Execution' column contains concrete values for variables x, y, and z. The 'Symbolic Execution' column contains symbolic representations of these variables. The 'path condition' column contains the path condition derived from the symbolic state.

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 22	x = x_0	$2*y_0 \neq x_0$
y = 7	y = y_0	
z = 14	z = $2*y_0$	

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 22$	$x = x_0$	$2*y_0 \neq x_0$
$y = 7$	$y = y_0$	
$z = 14$	$z = 2*y_0$	

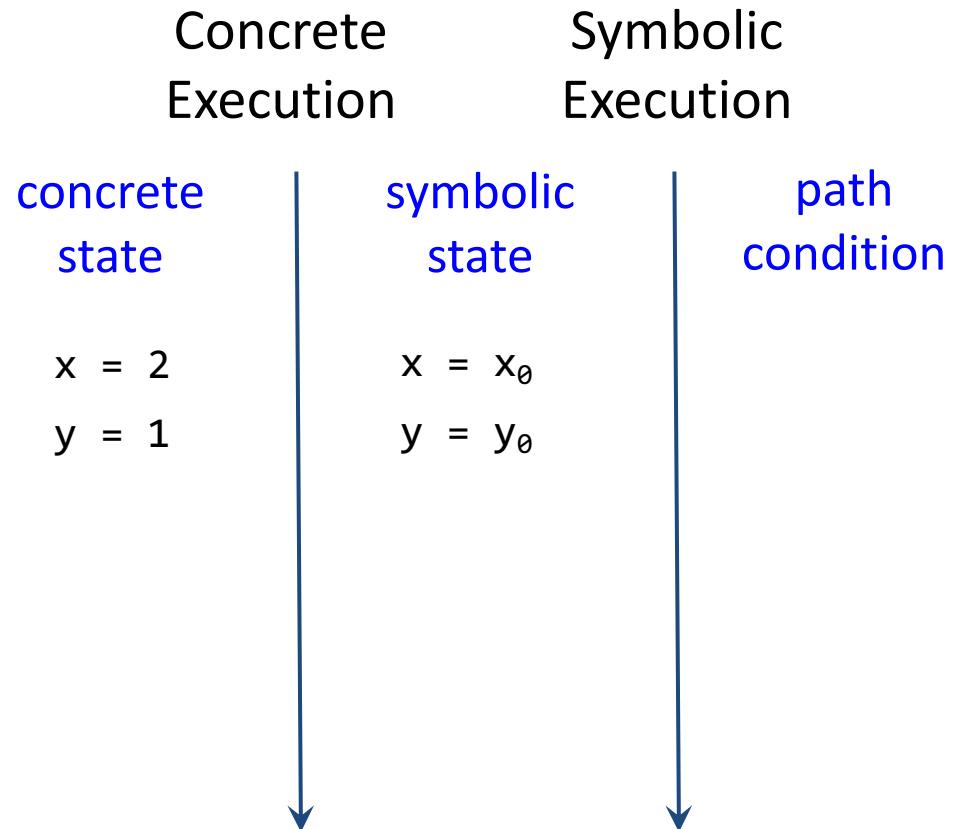
Solve: $2*y_0 == x_0$

Solution: $x_0 = 2, y_0 = 1$

The diagram illustrates the flow of information during symbolic execution. It starts with concrete execution on the left, which provides initial values for variables (x=22, y=7). These values are mapped to symbolic states (x=x₀, y=y₀) in the middle column. A path condition is derived from these symbolic states: 2*y₀ ≠ x₀. This path condition is then solved to find specific values for x₀ and y₀, resulting in the solution x₀ = 2, y₀ = 1.

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y); ←  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 2$	$x = x_0$	
$y = 1$	$y = y_0$	
$z = 2$	$z = 2*y_0$	

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 2$	$x = x_0$	$2*y_0 == x_0$
$y = 1$	$y = y_0$	
$z = 2$	$z = 2*y_0$	

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 2$	$x = x_0$	$2*y_0 == x_0$
$y = 1$	$y = y_0$	
$z = 2$	$z = 2*y_0$	$x_0 \leq y_0 + 10$

An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 2$	$x = x_0$	$2*y_0 == x_0$
$y = 1$	$y = y_0$	
$z = 2$	$z = 2*y_0$	$x_0 \leq y_0 + 10$

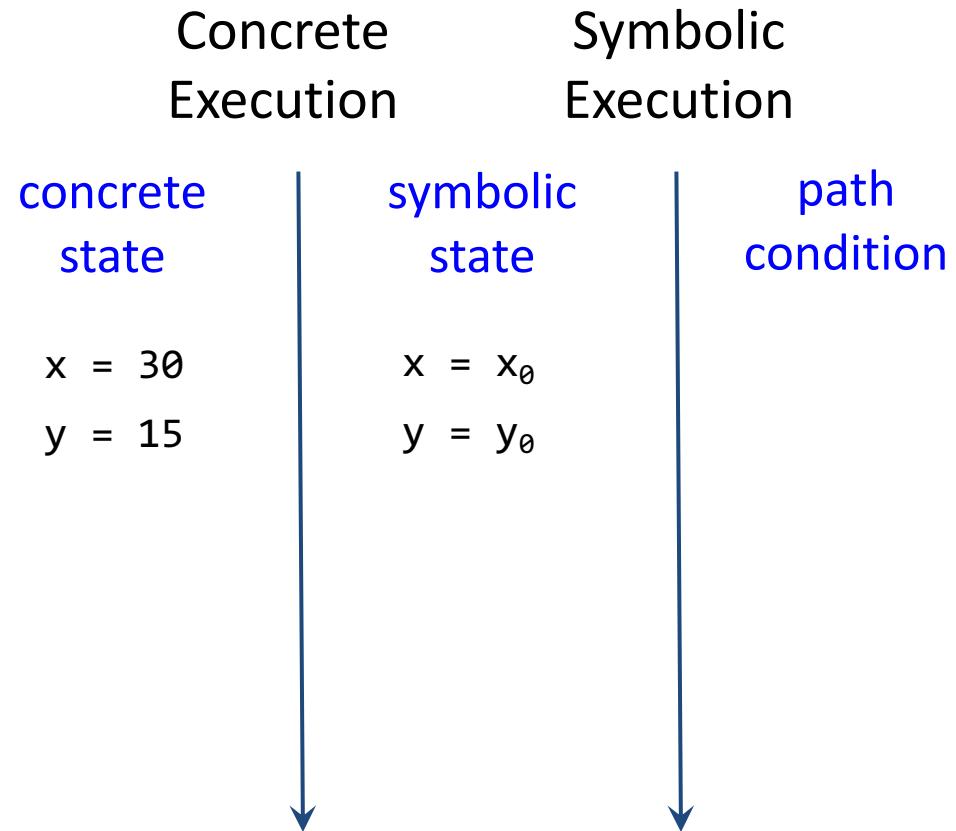
Solve: $(2*y_0 == x_0)$ and $(x_0 > y_0 + 10)$

Solution: $x_0 = 30$, $y_0 = 15$



An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y); ←  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



An Illustrative Example

```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y); ←  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 30	x = x_0	
y = 15	y = y_0	
z = 30	z = $2*y_0$	

An Illustrative Example

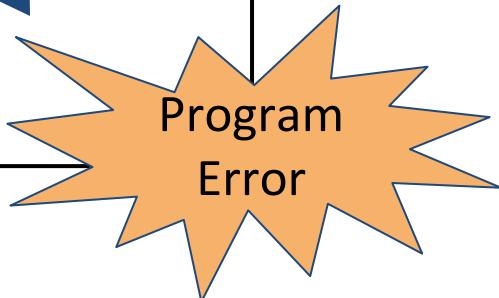
```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 30	x = x_0	$2*y_0 == x_0$
y = 15	y = y_0	
z = 30	z = $2*y_0$	

An Illustrative Example

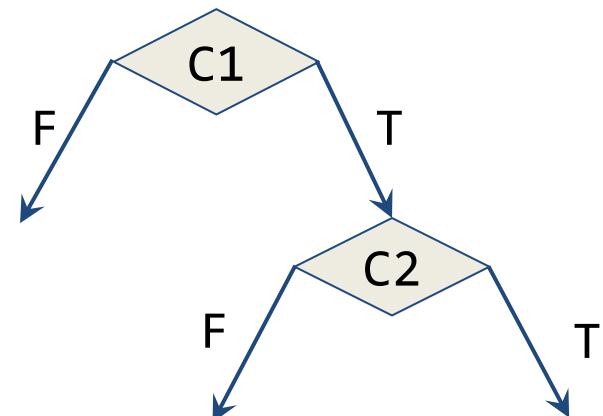
```
int foo(int v) {  
    return 2*v;  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 30	x = x_0	$2*y_0 == x_0$
y = 15	y = y_0	
z = 30	z = $2*y_0$	$x_0 > y_0 + 10$



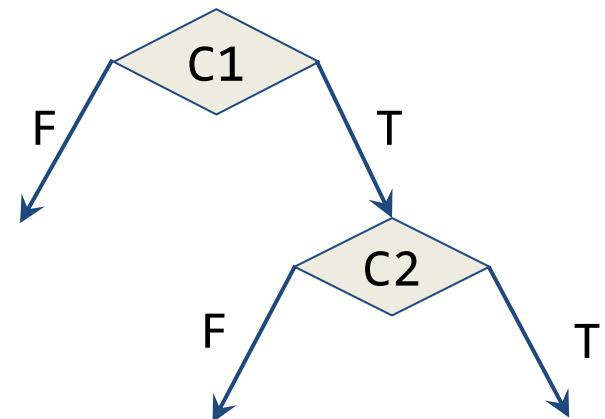
QUIZ: Computation Tree

Check all constraints that DSE might possibly solve in exploring the computation tree shown below:

 C1 $C1 \wedge C2$ C2 $C1 \wedge \neg C2$ $\neg C1$ $\neg C1 \wedge C2$ $\neg C2$ $\neg C1 \wedge \neg C2$ 

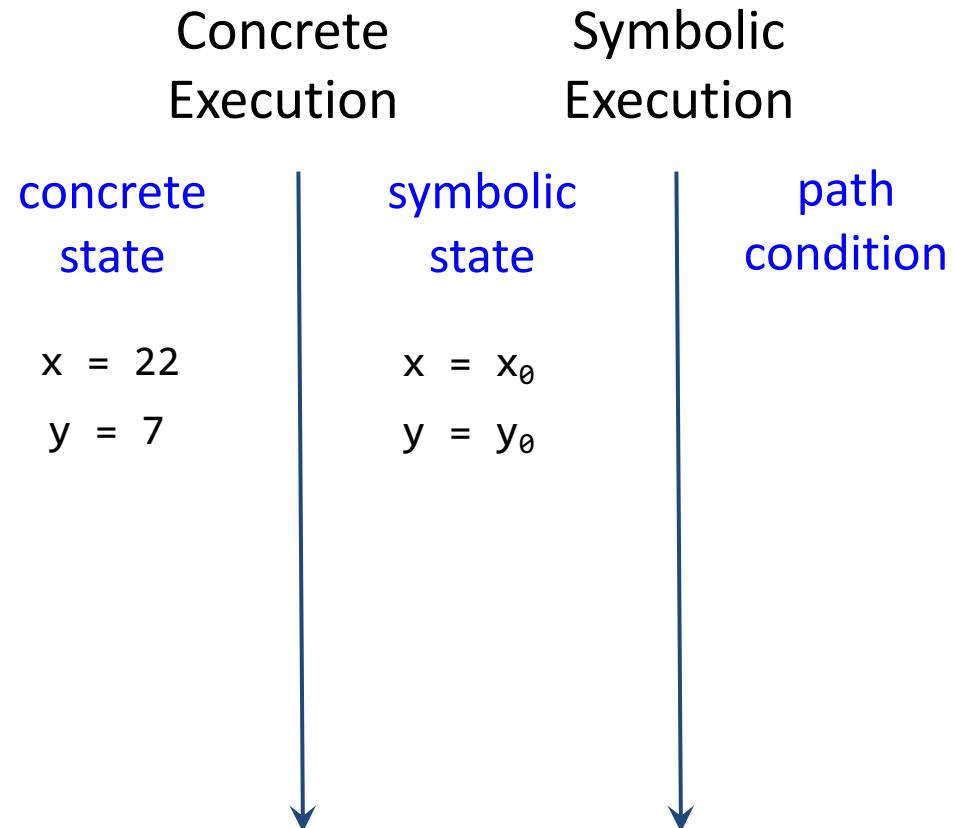
QUIZ: Computation Tree

Check all constraints that DSE might possibly solve in exploring the computation tree shown below:

 C1 $C1 \wedge C2$ C2 $C1 \wedge \neg C2$ $\neg C1$ $\neg C1 \wedge C2$ $\neg C2$ $\neg C1 \wedge \neg C2$ 

A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y); ←  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 22	x = x_0	
y = 7	y = y_0	
z = 601...129	z = $\text{secure_hash}(y_0)$	

The diagram illustrates the flow of data from concrete execution to symbolic execution. It consists of three columns: 'Concrete Execution' (left), 'Symbolic Execution' (middle), and 'path condition' (right). Vertical arrows connect the columns. The 'Concrete Execution' column contains variable assignments: x = 22, y = 7, and z = 601...129. The 'Symbolic Execution' column contains variable assignments: x = x_0 , y = y_0 , and z = $\text{secure_hash}(y_0)$. The 'path condition' column is currently empty. A blue arrow points from the 'Concrete Execution' column to the 'Symbolic Execution' column. Another blue arrow points from the 'Symbolic Execution' column to the 'path condition' column.

A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 22	x = x_0	$\text{secure_hash}(y_0)$
y = 7	y = y_0	$\neq x_0$
z = 601...129	z = $\text{secure_hash}(y_0)$	

Solve: $\text{secure_hash}(y_0) == x_0$

Don't know how to solve! Stuck?

A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)
```

Not stuck! Use
concrete state: replace
 y_0 by 7

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
$x = 22$	$x = x_0$	$secure_hash(y_0)$
$y = 7$	$y = y_0$	$\neq x_0$
$z = 601\dots129$	$z = secure_hash(y_0)$	

Solve: $secure_hash(y_0) == x_0$

Don't know how to solve! Stuck?

A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 22	x = x_0	secure_hash(y_0)
y = 7	y = y_0	$\neq x_0$
z = 601...129	z = secure_hash(y_0)	

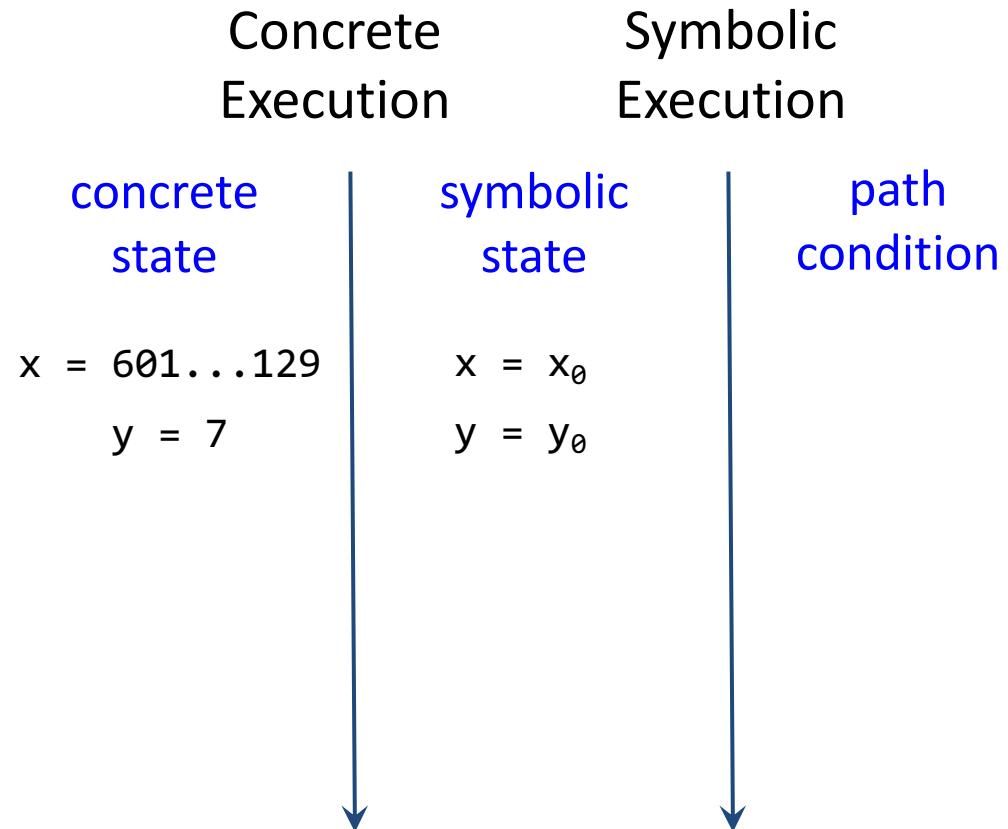
Solve: 601...129 == x_0

Solution: $x_0 = 601\dots129$, $y_0 = 7$

The diagram illustrates the flow of information during symbolic execution. It starts with concrete execution on the left, which provides initial values for variables x, y, and z. These values are then mapped to symbolic states on the right, where x is represented as x_0 and y as y_0 . The path condition is derived from the requirement that the result of the secure hash function must equal x. This leads to a solve step in a yellow box, where the equation 601...129 == x_0 is solved, resulting in the solution $x_0 = 601\dots129$ and $y_0 = 7$. Arrows indicate the flow from concrete to symbolic execution, and from symbolic execution to the solve step.

A More Complex Example

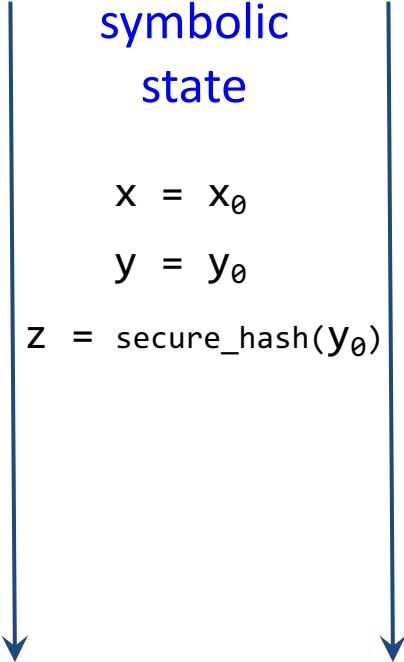
```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y); ←  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 601...129	x = x_0	
y = 7	y = y_0	
z = 601...129	z = $\text{secure_hash}(y_0)$	



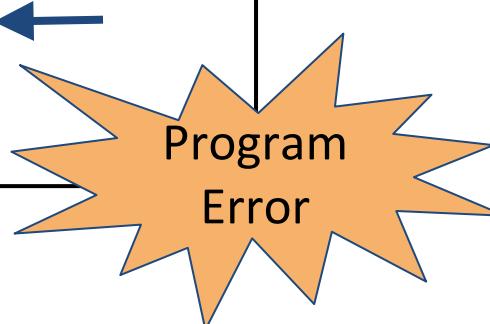
A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 601...129	x = x_0	secure_hash(y_0)
y = 7	y = y_0	$\equiv x_0$
z = 601...129	z = $\text{secure_hash}(y_0)$	

A More Complex Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    int z = foo(y);  
    if (z == x)  
        if (x > y+10)  
            ERROR;  
}
```



Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 601\dots129$	$x = x_0$	$\text{secure_hash}(y_0)$
$y = 7$	$y = y_0$	$\equiv x_0$
$z = 601\dots129$	$z = \text{secure_hash}(y_0)$	$x_0 > y_0 + 10$

QUIZ: Example Application

DSE tests the below program starting with input $x = 1$.

What is the input and constraint $(C1 \wedge C2 \wedge C3)$ solved in each run of DSE? Use depth-first search and leave trailing constraints blank if unused.

Run	x	C1	C2	C3
1	1	$5 \neq x_0$	$7 \neq x_0$	$9 == x_0$
2				
3				
4				

```
int test_me(int x) {
    int[] A = { 5, 7, 9 };
    int i = 0;
    while (i < 3) {
        if (A[i] == x) break;
        i++;
    }
    return i;
```

QUIZ: Example Application

DSE tests the below program starting with input $x = 1$.

What is the input and constraint $(C1 \wedge C2 \wedge C3)$ solved in each run of DSE? Use depth-first search and leave trailing constraints blank if unused.

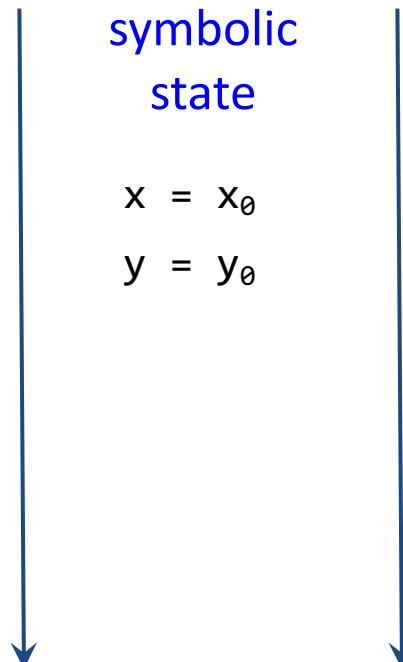
Run	x	C1	C2	C3
1	1	$5 \neq x_0$	$7 \neq x_0$	$9 == x_0$
2	9	$5 \neq x_0$	$7 == x_0$	
3	7	$5 == x_0$		
4	5			

```
int test_me(int x) {
    int[] A = { 5, 7, 9 };
    int i = 0;
    while (i < 3) {
        if (A[i] == x) break;
        i++;
    }
    return i;
}
```

A Third Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    if (x != y) ←  
        if (foo(x) == foo(y))  
            ERROR;  
}
```

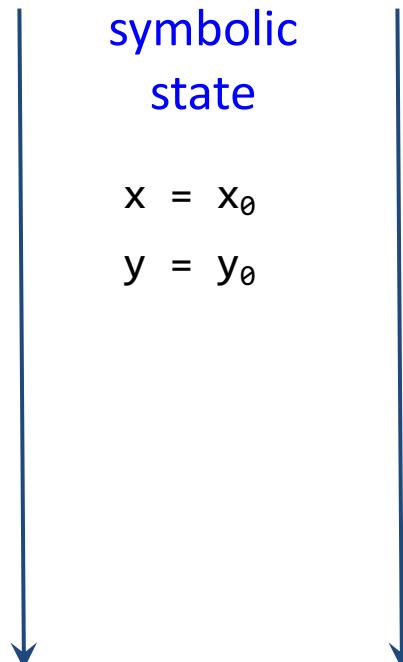
Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 22 y = 7	x = x_0 y = y_0	



A Third Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    if (x != y) ←  
        if (foo(x) == foo(y))  
            ERROR;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 22	x = x_0	$x_0 \neq y_0$
y = 7	y = y_0	



A Third Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    if (x != y)  
        if (foo(x) == foo(y))  
            ERROR;  
}
```



Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	x ₀ != y ₀
x = 22	x = x ₀	secure_hash(x ₀) !=
y = 7	y = y ₀	secure_hash(y ₀)

Solve: $x_0 \neq y_0$ and
 $\text{secure_hash}(x_0) == \text{secure_hash}(y_0)$

Use concrete state: replace y_0 by 7.

A Third Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    if (x != y)  
        if (foo(x) == foo(y))  
            ERROR;  
}
```



Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 22	x = x_0	$x_0 \neq y_0$
y = 7	y = y_0	secure_hash(x_0) != secure_hash(y_0)

Solve: $x_0 \neq 7$ and
secure_hash(x_0) == 601...129

Use concrete state: replace x_0 by 22.

A Third Example

```
int foo(int v) {  
    return secure_hash(v);  
}  
  
void test_me(int x, int y) {  
    if (x != y)  
        if (foo(x) == foo(y))  
            ERROR;  
}
```

False negative!

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
x = 22	x = x_0	$x_0 \neq y_0$
y = 7	y = y_0	secure_hash(x_0) != secure_hash(y_0)

Solve: $22 \neq 7$ and
 $438\dots861 == 601\dots129$

Unsatisfiable!

QUIZ: Properties of DSE

Assume that programs can have infinite computation trees.
Which statements are true of DSE applied to such programs?

- DSE is guaranteed to terminate.
- DSE is complete: if it ever reaches an error, the program can reach that error in some execution.
- DSE is sound: if it terminates and did not reach an error, the program cannot reach an error in any execution.

QUIZ: Properties of DSE

Assume that programs can have infinite computation trees.
Which statements are true of DSE applied to such programs?

- DSE is guaranteed to terminate.
- DSE is complete: if it ever reaches an error, the program can reach that error in some execution.
- DSE is sound: if it terminates and did not reach an error, the program cannot reach an error in any execution.

Another Example: Testing Data Structures

- Random Test Driver:
 - random value for `x`
 - random memory graph reachable from `p`
- Probability of reaching **ERROR** is extremely low

```
typedef struct cell {  
    int data;  
    struct cell *next;  
} cell;  
  
int foo(int v) { return 2*v + 1; }  
  
int test_me(int x, cell *p) {  
    if (x > 0)  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
    return 0;  
}
```

Data-Structure Example

```
typedef struct cell {  
    int data;  
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} cell;  
  
int foo(int v) { return 2*v + 1; }  
  
int test_me(int x, cell *p) {  
    if (x > 0) ←  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 236 p = NULL	x = x_0 p = p_0	

Data-Structure Example

```
typedef struct cell {  
    int data;  
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} cell;  
  
int foo(int v) { return 2*v + 1; }  
  
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    if (x > 0) ←  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 236 p = NULL	x = x_0 p = p_0	
		$x_0 > 0$

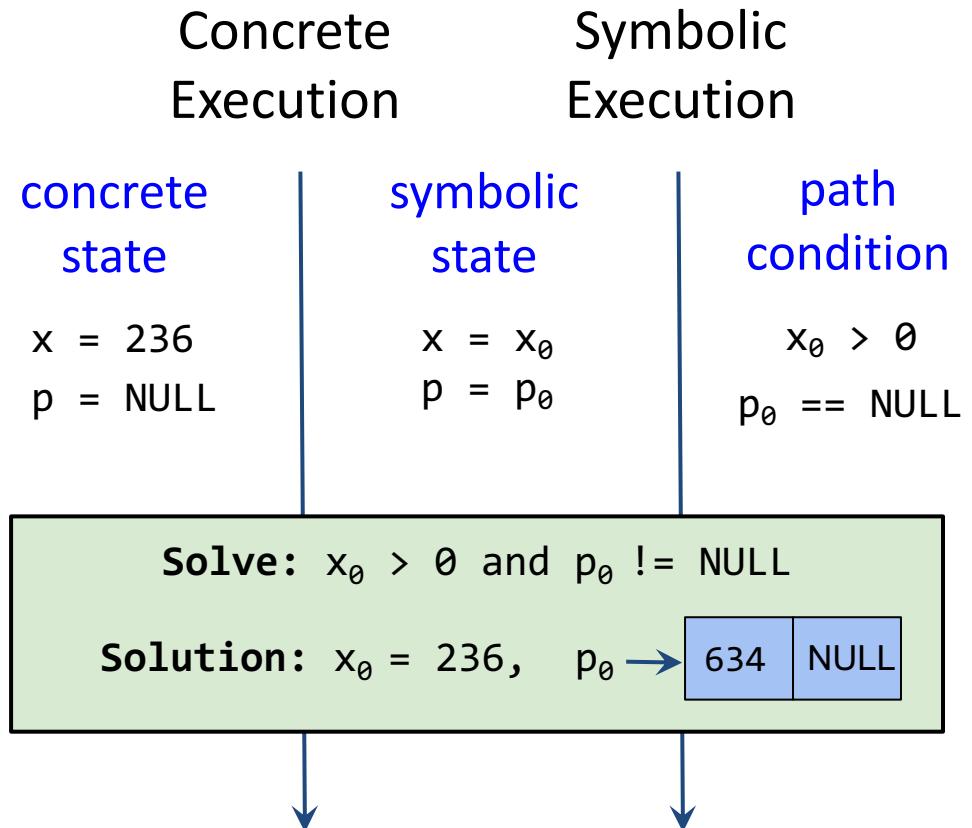
Data-Structure Example

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} cell;  
  
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    if (x > 0)  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0; ←  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
x = 236 p = NULL	x = x_0 p = p_0	$x_0 > 0$ $p_0 == \text{NULL}$

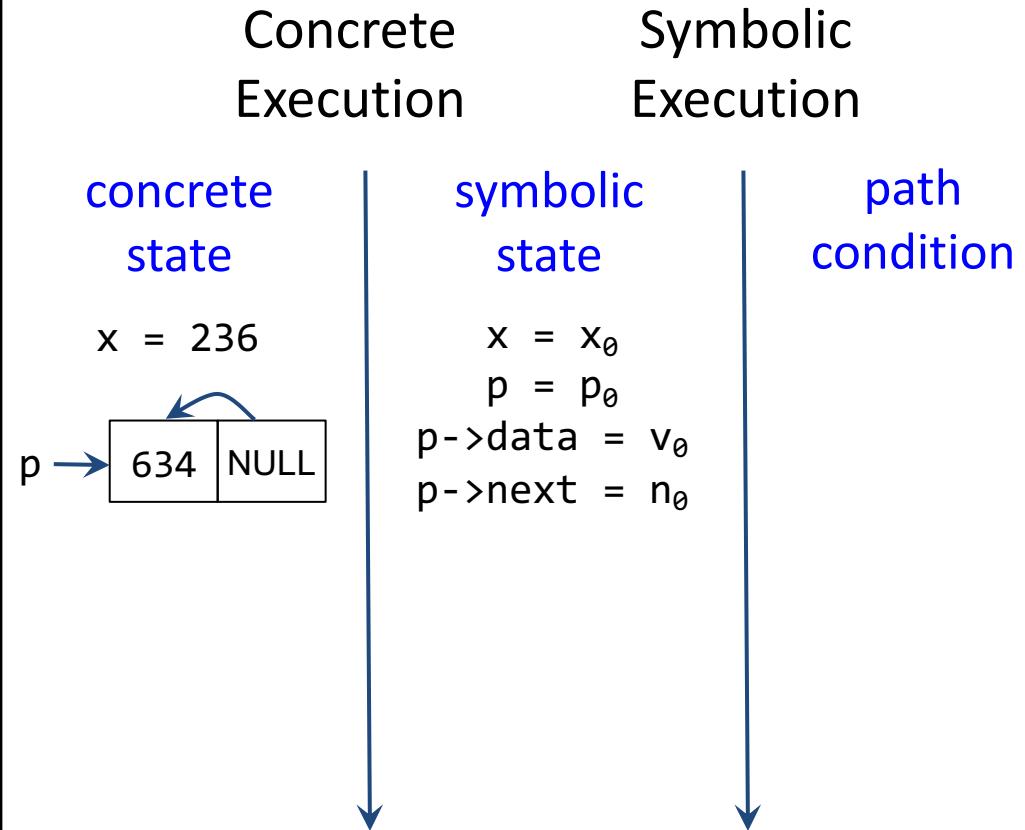
Data-Structure Example

```
typedef struct cell {  
    int data;  
    struct cell *next;  
} cell;  
  
int foo(int v) { return 2*v + 1; }  
  
int test_me(int x, cell *p) {  
    if (x > 0)  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0; ←  
}
```



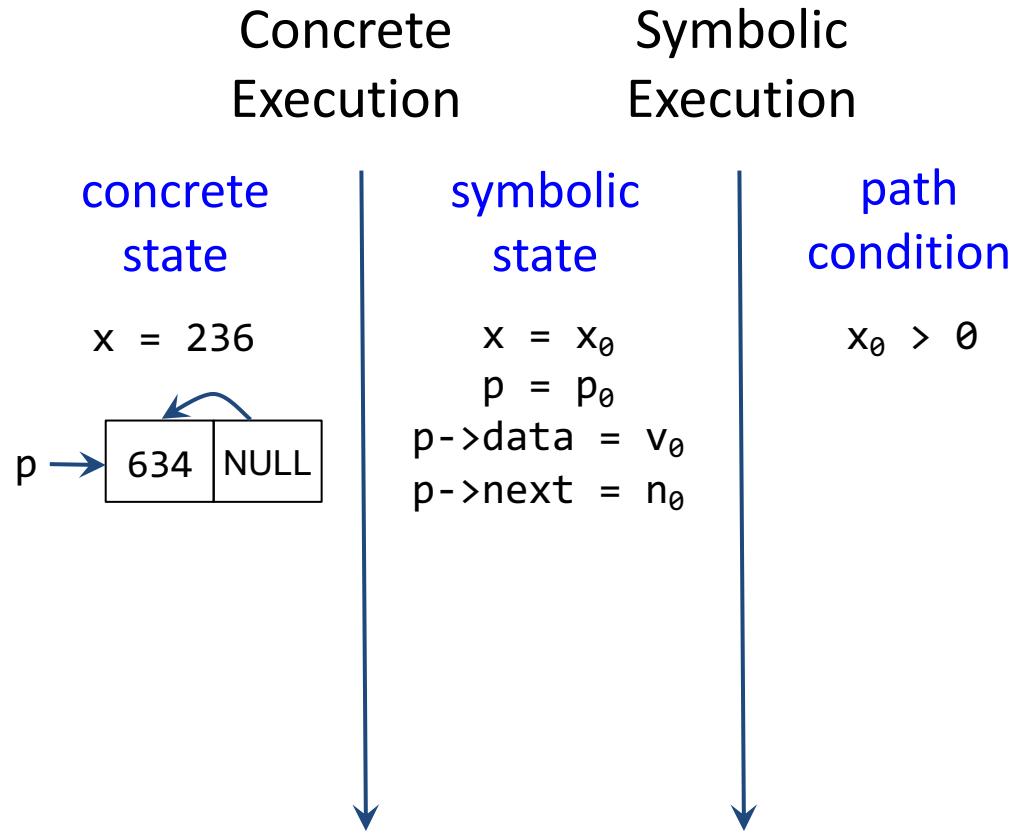
Data-Structure Example

```
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} cell;  
  
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    if (x > 0) ←  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```



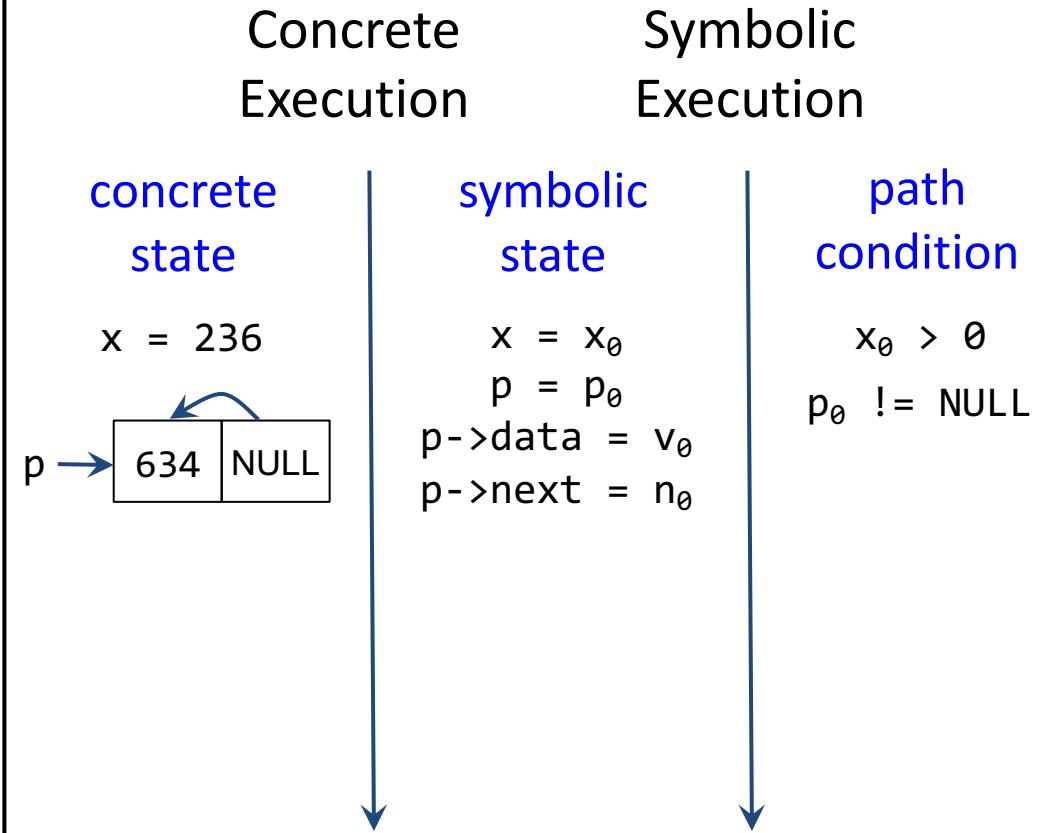
Data-Structure Example

```
typedef struct cell {  
    int data;  
    struct cell *next;  
} cell;  
  
int foo(int v) { return 2*v + 1; }  
  
int test_me(int x, cell *p) {  
    if (x > 0) ←  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```



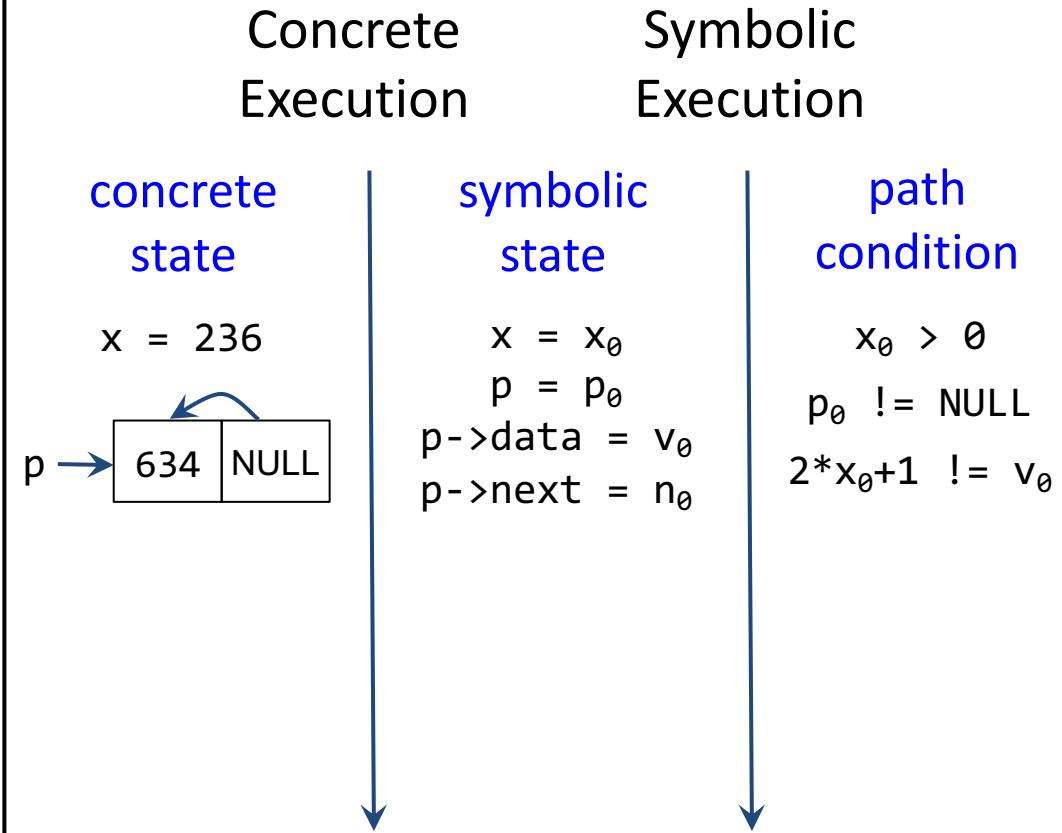
Data-Structure Example

```
typedef struct cell {  
    int data;  
    struct cell *next;  
} cell;  
  
int foo(int v) { return 2*v + 1; }  
  
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    if (x > 0)  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```



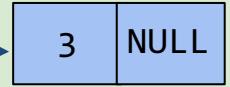
Data-Structure Example

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            if (foo(x) == p->data)  
                if (p->next == p)  
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    return 0; ←  
}
```



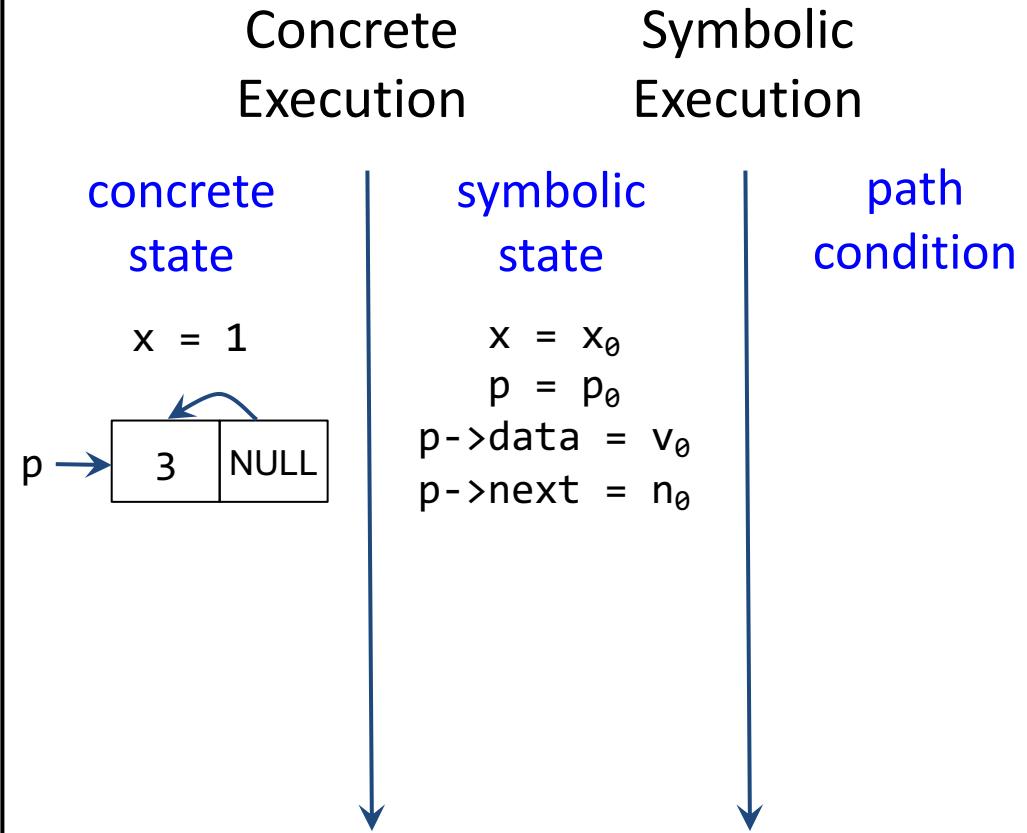
Data-Structure Example

```
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    struct cell *next;  
} cell;  
  
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                if (p->next == p)  
                    ERROR;  
  
    return 0; ←  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
$x = 236$ 	$x = x_0$ $p = p_0$ $p \rightarrow \text{data} = v_0$ $p \rightarrow \text{next} = n_0$	$x_0 > 0$ $p_0 \neq \text{NULL}$ $2*x_0+1 \neq v_0$
Solve: $x_0 > 0$ and $p_0 \neq \text{NULL}$ and $2*x_0+1==v_0$		
Solution: $x_0 = 1$, $p_0 \rightarrow$		

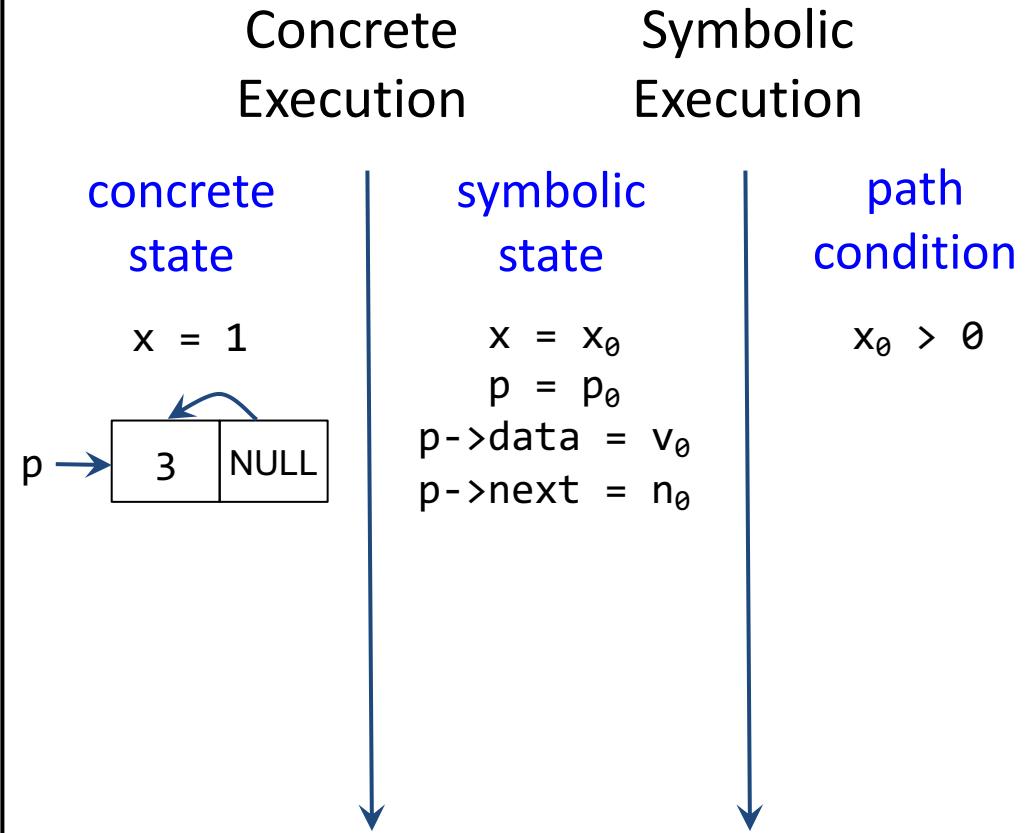
Data-Structure Example

```
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} cell;  
  
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int test_me(int x, cell *p) {  
    if (x > 0) ←  
        if (p != NULL)  
            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```



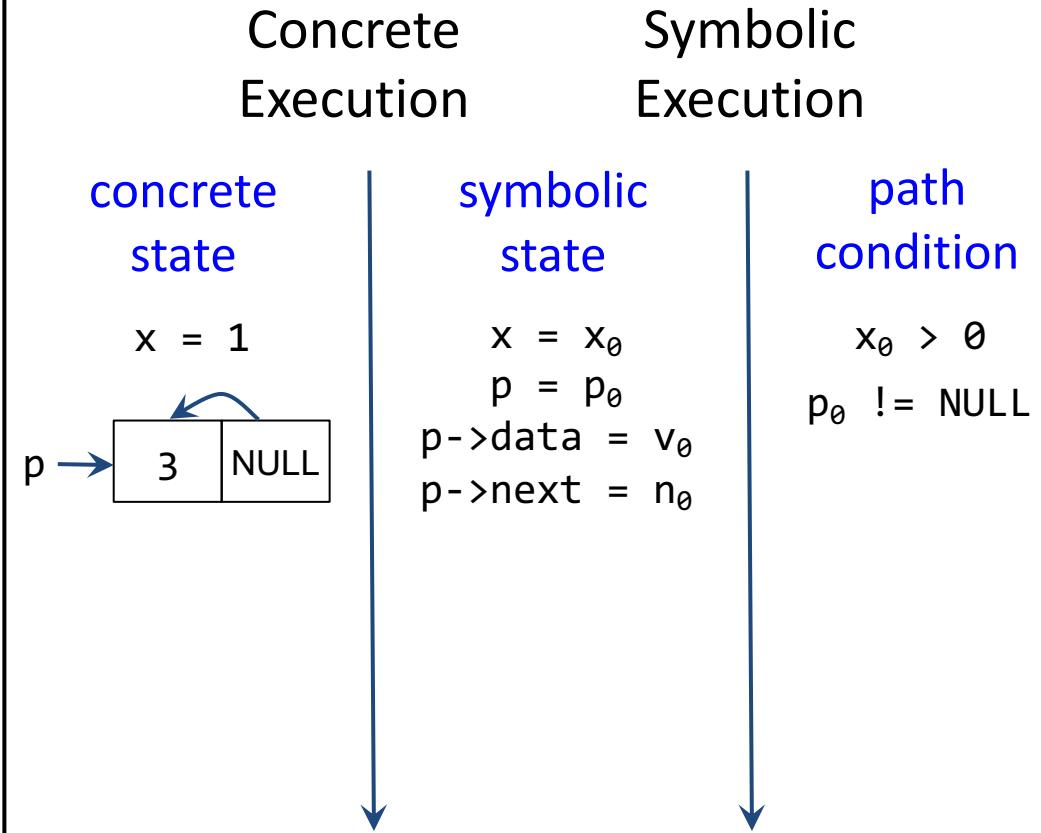
Data-Structure Example

```
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Data-Structure Example

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            if (foo(x) == p->data)  
                if (p->next == p)  
                    ERROR;  
  
    return 0;  
}
```



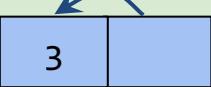
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```

Concrete Execution	Symbolic Execution	path condition
concrete state $x = 1$  $p \rightarrow \boxed{3 \text{ } \text{ } \text{NULL}}$	symbolic state $x = x_0$ $p = p_0$ $p \rightarrow \text{data} = v_0$ $p \rightarrow \text{next} = n_0$	$x_0 > 0$ $p_0 \neq \text{NULL}$ $2*x_0+1 == v_0$

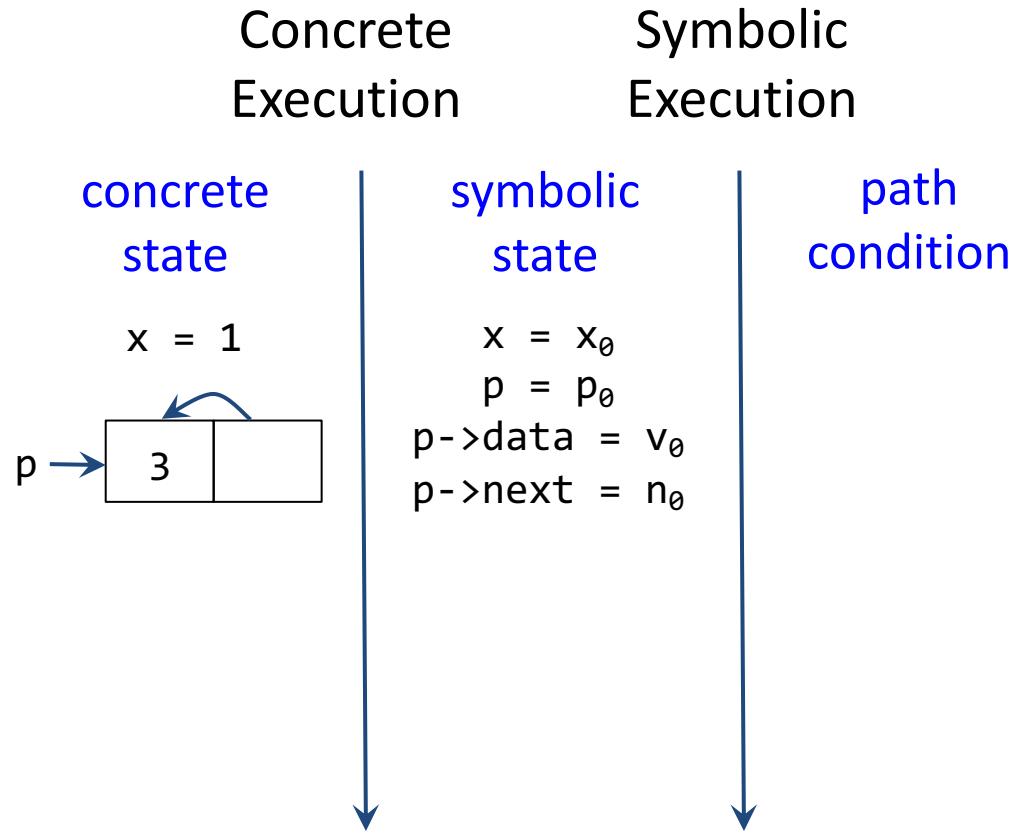
Data-Structure Example

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Concrete Execution	Symbolic Execution	path condition
concrete state	symbolic state	
$x = 1$	$x = x_0$	$x_0 > 0$
	$p = p_0$	$p_0 \neq \text{NULL}$
	$p->\text{data} = v_0$	$2*x_0+1 == v_0$
	$p->\text{next} = n_0$	$n_0 \neq p_0$
Solve: $x_0 > 0$ and $p_0 \neq \text{NULL}$ and $2*x_0+1==v_0$ and $n_0 == p_0$		
Solution: $x_0 = 1$, $p_0 \rightarrow$ 		

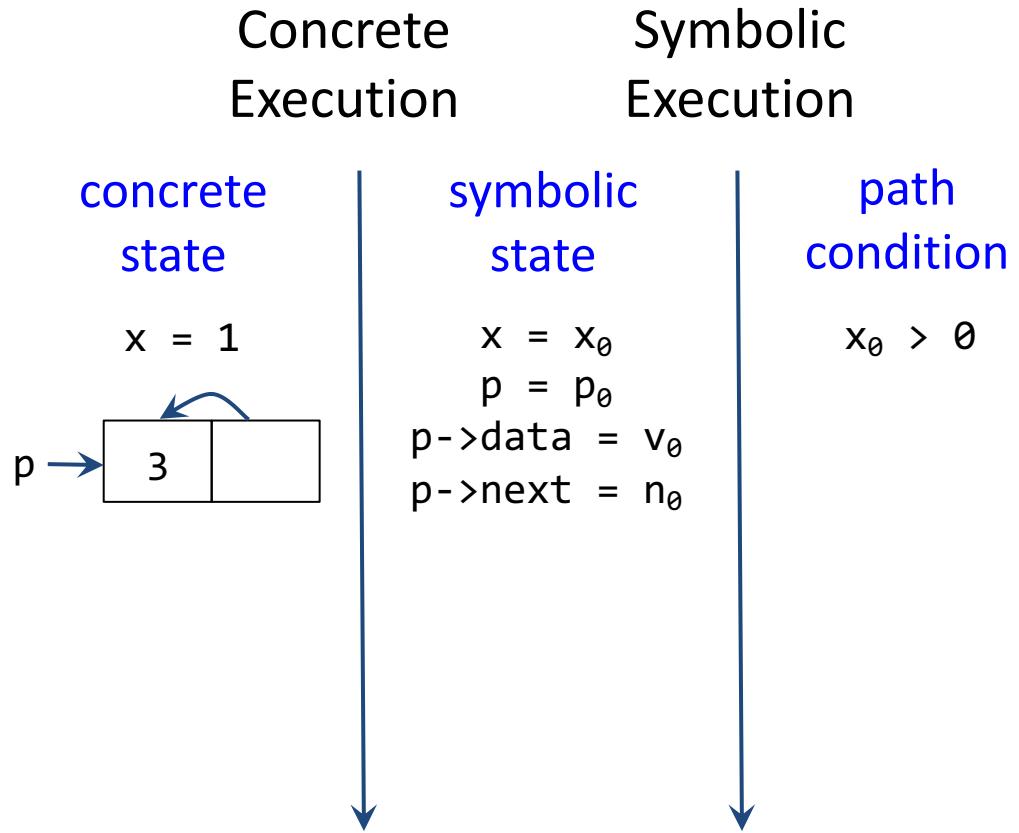
Data-Structure Example

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```



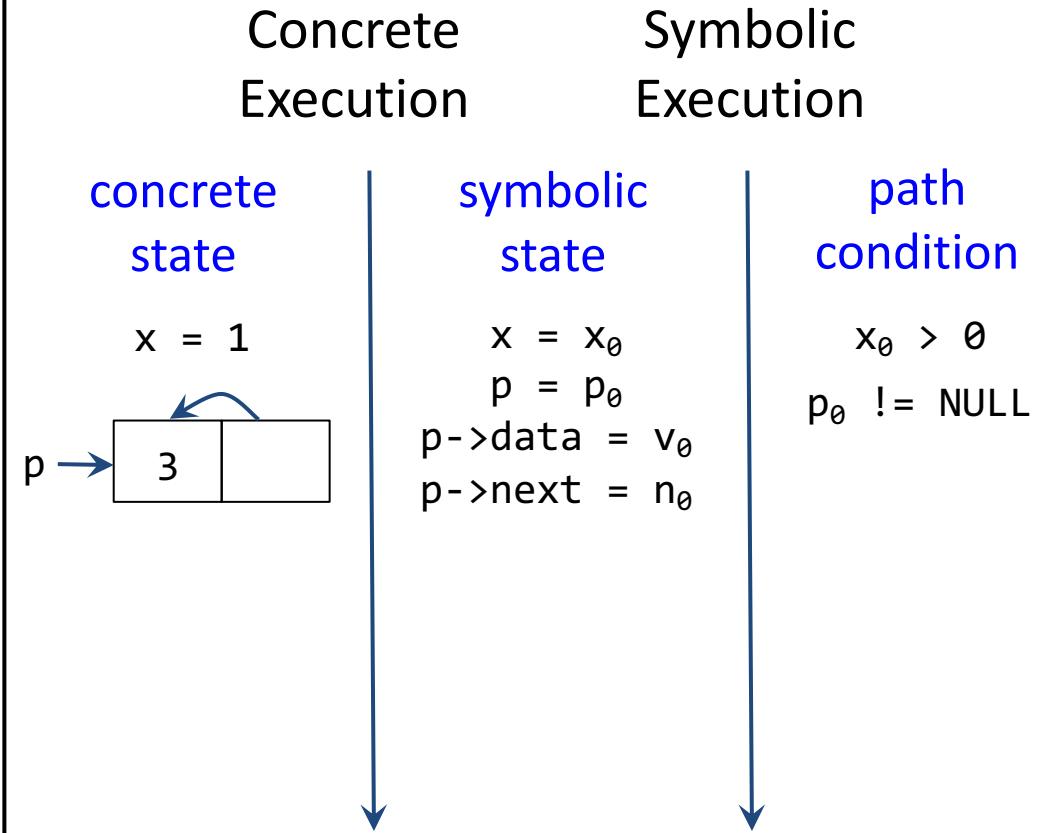
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```



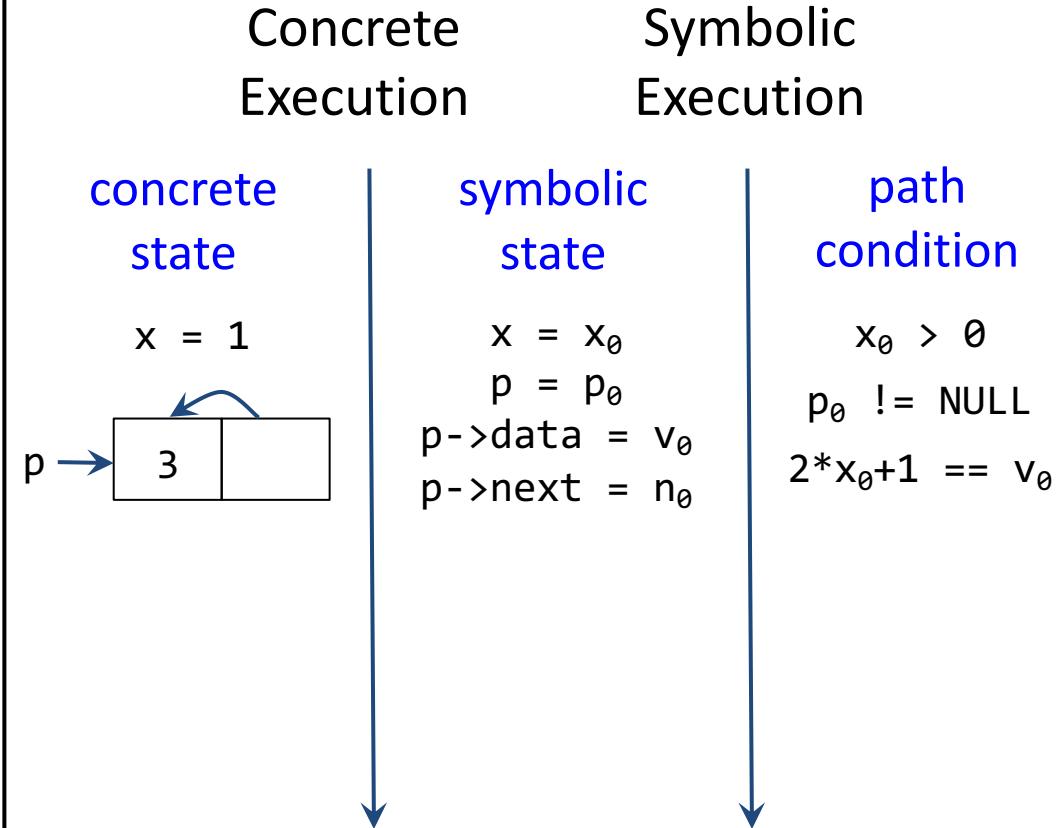
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Data-Structure Example

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                    ERROR;  
    return 0;  
}
```

Concrete Execution	Symbolic Execution	path condition
concrete state $x = 1$ 	symbolic state $x = x_0$ $p = p_0$ $p->data = v_0$ $p->next = n_0$	$x_0 > 0$ $p_0 \neq \text{NULL}$ $2*x_0+1 == v_0$ $n_0 == p_0$



Approach in a Nutshell

- Generate concrete inputs, each taking different program path
- On each input, execute program both **concretely** and **symbolically**
- Both **cooperate** with each other:
 - Concrete execution **guides** symbolic execution
 - Enables it to overcome incompleteness of theorem prover
 - Symbolic execution **guides** generation of concrete inputs
 - Increases program code coverage

QUIZ: Characteristics of DSE

- The testing approach of DSE is:
 - Automated, black-box
 - Automated, white-box
 - Manual, black-box
 - Manual, white-box
- The input search of DSE is:
 - Randomized
 - Systematic
- The static analysis of DSE is:
 - Flow-insensitive
 - Flow-sensitive
 - Path-sensitive
- The instrumentation in DSE is:
 - Sampled
 - Non-sampled

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Case Study: SGLIB C Library

- Found two bugs in sglib 1.0.1
 - reported to authors, fixed in sglib 1.0.2
- Bug 1: doubly-linked list
 - segmentation fault occurs when a non-zero length list is concatenated with zero-length list
 - discovered in 140 iterations (< 1 second)
- Bug 2: hash-table
 - an infinite loop in hash-table is_member function
 - 193 iterations (1 second)

Case Study: SGLIB C Library

Name	Run time (sec.)	# iterations	# branches explored	% branch coverage	# functions tested	# bugs found
Array Quick Sort	2	732	43	97.73	2	0
Array Heap Sort	4	1764	36	100.00	2	0
Linked List	2	570	100	96.15	12	0
Sorted List	2	1020	110	96.49	11	0
Doubly Linked List	3	1317	224	99.12	17	1
Hash Table	1	193	46	85.19	8	1
Red Black Tree	2629	1,000,000	242	71.18	17	0

Case Study: Needham-Schroeder Protocol

- Tested a C implementation of a security protocol (Needham-Schroeder) with a known (man-in-the-middle) attack
 - 600 lines of code
 - Took fewer than 13 seconds on a machine with 1.8 GHz processor and 2 GB RAM to discover the attack
- In contrast, a software model-checker (VeriSoft) took 8 hours

Realistic Implementations

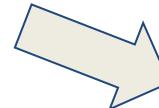
- **KLEE**: LLVM (C family of languages)
- **PEX**: .NET Framework
- **jCUTE**: Java
- **Jalangi**: Javascript
- **SAGE** and **S2E**: binaries (x86, ARM, ...)

Case Study: SAGE Tool at Microsoft

- SAGE = Scalable Automated Guided Execution
- Found many expensive security bugs in many Microsoft applications (*Windows, Office, etc.*)
- Used daily in various Microsoft groups, runs 24/7 on 100's of machines
- What makes it so useful?
 - Works on *large applications* => finds bugs across components
 - Focus on input *file fuzzing* => fully automated
 - Works on *x86 binaries* => easy to deploy (not dependent on language or build process)

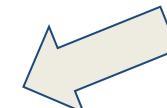
Example: SAGE Crashing a Media Parser

```
00000000h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000060h: 00 00 00 00 ; ....
```



```
00000000h: 52 49 46 46 00 00 00 00 00 00 00 00 00 00 00 00 ; RIFF.....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000060h: 00 00 00 00 ; ....
```

```
00000000h: 52 49 46 46 00 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF...** .....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000060h: 00 00 00 00 ; ....
```



... after a few more iterations:

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 00 ; RIFF=...*** .....  
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; ....  
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 76 69 64 73 ; ....strh....vids  
00000040h: 00 00 00 00 73 74 72 66 B2 75 76 3A 28 00 00 00 ; ....strf^uv:(...  
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 00 ; ....  
00000060h: 00 00 00 00 ; ....
```

What Have We Learned?

- What is (dynamic) symbolic execution?
- Systematically generate (numeric and pointer) inputs
- Computation tree and error reachability
- Tracking concrete state, symbolic state, path condition
- Combined dynamic and static analysis =>
Hybrid analysis
- Complete, but no soundness or termination
guarantees

Paper Readings

- *DART: Directed Automated Random Testing*. PLDI'05
<https://people.eecs.berkeley.edu/~ksen/papers/dart.pdf>
- *KLEE: Unassisted and Automatic Generation of High-Coverage Tests for Complex Systems Programs*. OSDI'08
<https://llvm.org/pubs/2008-12-OSDI-KLEE.pdf>
- Symbolic Execution for Software Testing: Three Decades Later (ACM'13)
<https://people.eecs.berkeley.edu/~ksen/papers/cacm13.pdf>